## Homework 2

Instructor: Xiangyu Guo
Deadline: 7/01/2020
Your Name: $\qquad$ Your Student ID: $\qquad$

| Problems | 1 | 2 | 3 | 4 | Total |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Max. Score | 10 | 20 | 20 | 20 | 70 |
| Your Score |  |  |  |  |  |

Problem 1 (10 points). Consider the following graph $G$ with non-negative edge weights. Use Dijkstra's algorithm to compute the shortest paths from $s$ to all other

vertices in $G$. Fill the following table to describe the execution of the algorithm. The algorithm maintains a set $\Gamma$ of vertices, of which vertices have their $d(\cdot)$ value equals the shortest path length. The $d$ value of a vertex $v \notin \Gamma$ is $\min _{u \in S:(u, v) \in E}(d(u)+w(u, v))$. The $\pi$ value of a vertex $v$ is the vertex $u \in \Gamma$ such that $d(v)=d(u)+w(u, v)$; if $d(v)=\infty$, then $\pi(v)=" / "$.

| iteration | vertex added to $\Gamma$ | $a$ |  | $b$ |  | c |  | $d$ |  | $e$ |  | $f$ |  | $g$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $d$ | $\pi$ | $d$ | $\pi$ | $d$ | $\pi$ | $d$ | $\pi$ | $d$ | $\pi$ | $d$ | $\pi$ | $d$ | $\pi$ |
| 1 | $s$ | $\infty$ | / | 9 | $s$ | 17 | $s$ | 2 | $s$ | $\infty$ | / | $\infty$ | / | $\infty$ | / |
| 2 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 3 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 4 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 5 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 6 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 7 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 8 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

Table 1: Dijkstra's algorithm for Shortest Path

Problem 2 (20 points). Assume we are given an undirected graph $G=(V, E)$ with non-negative edge weights $\left(w_{e}\right)_{e \in E}$, and two vertices $s$ and $t$ in $V$.
(2a) (10 points) Let $T$ be the unique minimum spanning tree of $G$. Is the following statement true or false? If we change the weight of every edge $e$ from $w_{e}$ to $w_{e}^{2}$, then $T$ is still the unique minimum spanning tree of $G$. Justify your answer.
(2b) (10 points) Let $P$ be the unique shortest path from $s$ to $t$. Is the following statement true or false? If we change the weight of every edge $e$ from $w_{e}$ to $w_{e}^{2}$, then $P$ is still the unique shortest path from $s$ to $t$. Justify your answer.

Problem 3 (20 points). Balanced strings are those who have equal quantity of "L" and "R" characters. Given a balanced string $s$, the goal is to split it into the maximum amount of balanced strings. For example, if $s=$ "RLRRLLRLRL", the optimal split is splitting into "RL", "RRLL", "RL", "RL", each substring contains same number of "L" and "R". Another example is $s=$ "LRRLLRLR", where there're multiple ways to split $s$, but the only optimal way is "LR" + "RL"+"LR" +"LR".

- (5 points) Suppose your greedy strategy is to pick the first few characters from $s$. Which characters are you going to choose?
- (15 points) Prove the safety property of your greedy strategy. (Hint: note that $s$ itself is already balanced.)

Problem (20 points). Consider a long country road with houses scattered very sparsely along it. You may picture the road as a long straight line segment, with the starting point (mile stone 0 ) and the endpoint (mile stone $L$ ). Each house is identified by its distance to the western endpoint. A cell phone company wants to set up cell phone services along the road. The company can place a base station at any house. (The monthly charge will be waived if a base station is located in a house, so the house owners are eager to accommodate base stations). The power of base stations are limited that can only cover a distance of 5 miles. The goal for the company is to select a minimum number of base stations so that every house on the road is within 5 miles of a base station. (The bases stations are connected by other means, say by Satellite. So the distance between them can be more than 5 miles). See Figure 1 for a example.


Figure 1: Using 5 base stations to cover all houses in $A$ (denoted by the solid circles on the line).

A formal description of the problems: The input is an array $A$ of n points: $A=$ $a_{1}<a_{2}<\ldots<a_{n}$, where each $a_{i}(1 \leq i \leq n)$ represents a house. We need to select a subset $B \subseteq A$ such that: (1) for every point $a_{i} \in A$, there is a point $a_{j} \in B$ with
$\left|a_{i}-a_{j}\right| \leq 5$, and (2) the size of $B$ is minimum, subject to condition (1). Describe a greedy algorithm for solving this problem. You need to prove the correctness of the algorithm.

- (5 points) Suppose your greedy strategy picks base station locations from left to right. Where would you set up your first base station?
- (15 points) Prove the safety property of your greedy strategy.

