

## Homework 2

Recall:

**Definition 1.** A category  $\mathcal{C}$  is *complete* (respectively *cocomplete*) if every functor  $F : \mathcal{D} \rightarrow \mathcal{C}$  from a small category  $\mathcal{D}$  has a limit (respectively colimit). ◆

On the other hand, we can require less:

**Definition 2.** A category  $\mathcal{C}$  *has (co)products* if arbitrary (small) families of objects in  $\mathcal{C}$  admit a (co)product in  $\mathcal{C}$ . Equivalently, every functor  $F : \mathcal{D} \rightarrow \mathcal{C}$  for discrete small  $\mathcal{D}$  admits a (co)limit.

Similarly,  $\mathcal{C}$  *has (co)equalizers* if every functor  $F : \mathcal{D} \rightarrow \mathcal{C}$  from the category

$$\mathcal{D} = \bullet \rightrightarrows \bullet$$

has a (co)limit. ◆

**Problem 1.** Show that a category is (co)complete if and only if it has (co)products and (co)equalizers.

Feel free to prove only the product/equalizer version, as the other result is entirely parallel by simply reversing the errors. One possible course of action I recommend:

Start with an arbitrary functor  $F : \mathcal{D} \rightarrow \mathcal{C}$  from a small category  $\mathcal{D}$  and form the products (which you know exist by the hypothesis)

$$P_o := \prod_{\text{objects } d \in \mathcal{D}} Fd$$

and

$$P_a := \prod_{\text{arrows } f: d \rightarrow d' \text{ in } \mathcal{D}} x_f \text{ where } x_f = Fd'$$

Then consider the two parallel morphisms

$$P_o \begin{array}{c} \xrightarrow{\psi_l} \\ \xrightarrow{\psi_r} \end{array} P_a$$

defined as follows :

For  $\psi_r$  and an arrow  $f : d \rightarrow d'$  in  $\mathcal{D}$ , the component  $P_o \rightarrow x_f$  of the morphism into the product  $P_a = \prod_f x_f$  is the composition

$$P_o \xrightarrow{\pi_{d'}} Fd' \xrightarrow{\text{id}} x_f = Fd'$$

where

$$\pi_{d'} : P_o = \prod_d Fd \rightarrow Fd'$$

is the structure morphism coming from the product structure.

On the other hand, for  $\psi_l$  and an arrow  $f : d \rightarrow d'$  the component  $P_o \rightarrow x_f$  is the composition

$$P_o \xrightarrow{\pi_d} Fd \xrightarrow{Ff} x_f = Fd'.$$

By assumption, the morphisms  $\psi_l$  and  $\psi_r$  have an equalizer

$$\bullet \xrightarrow{\iota} P_o \begin{array}{c} \xrightarrow{\psi_l} \\ \xrightarrow{\psi_r} \end{array} P_a$$

by assumption. Try to show that the morphisms

$$\bullet \xrightarrow{\iota} P_o \xrightarrow{\pi_d} Fd$$

for  $d \in \mathcal{D}$  constitute a cone from  $\bullet$  to  $F$ , and prove that this cone is universal (i.e. a limit for  $F$ ).

As a consequence of [Problem 1](#), prove:

**Problem 2.** *Let  $F : \mathcal{C} \rightarrow \mathcal{C}'$  be a functor between complete categories. Show that  $F$  is continuous if and only if it preserves products and equalizers.*

Here's a remark I made in passing at the end of the Feb 6 lecture, which one of you pointed out we hadn't gone over before; I'm making it into a problem:

**Problem 3.** *Let  $R$  and  $S$  be two rings and  $F : {}_R\text{Mod} \rightarrow {}_S\text{Mod}$  an additive functor. Prove that  $F$  is left exact if and only if it preserves equalizers (i.e. turns equalizers into equalizers).*

As an immediate consequence, continuous functors (such as right adjoints, say) are left exact.

**Problem 4.** *5.17 from Rotman.*

**Problem 5.** *5.32 from Rotman.*

**Problem 6.** *5.34 from Rotman.*